TOWARDS A CHARACTERIZATION OF THE COVERT CAPACITY OF BOSONIC CHANNELS UNDER TRACE DISTANCE

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COVERT COMMUNICATIONS

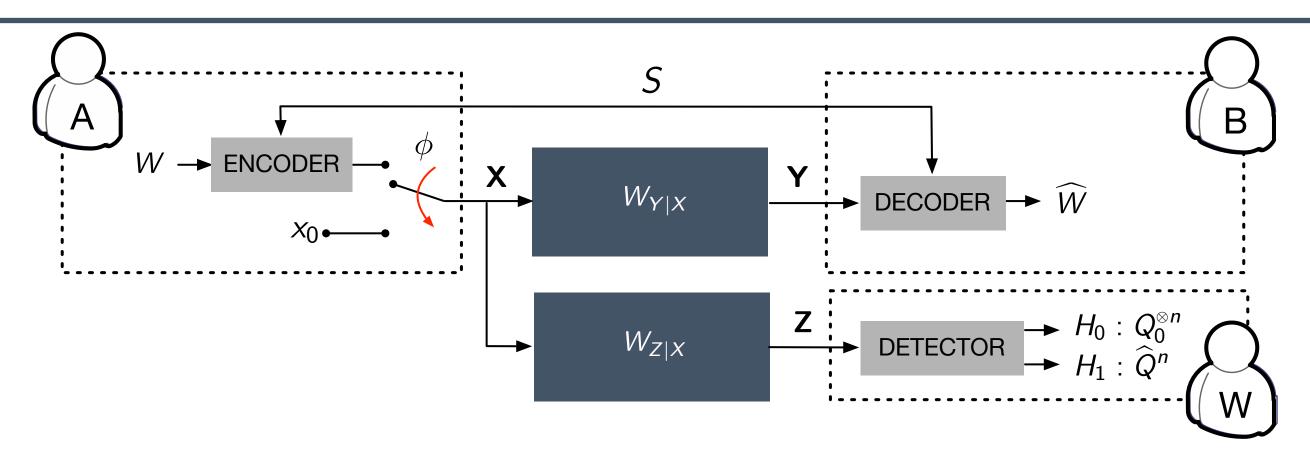
- Enhance privacy by avoiding exposure of information transmission
- Communication with low probability of detection

FUNDAMENTAL LIMITS OF COVERT COMMUNICATIONS

- (Square Root Law [Bash et al.'13]) No more than $\mathcal{O}(\sqrt{n})$ bits over n uses of noisy memoryless channel
- Zero-rate regime: rate of communication vanishes asymptotically

EXTENSIONS AND RELATED WORKS

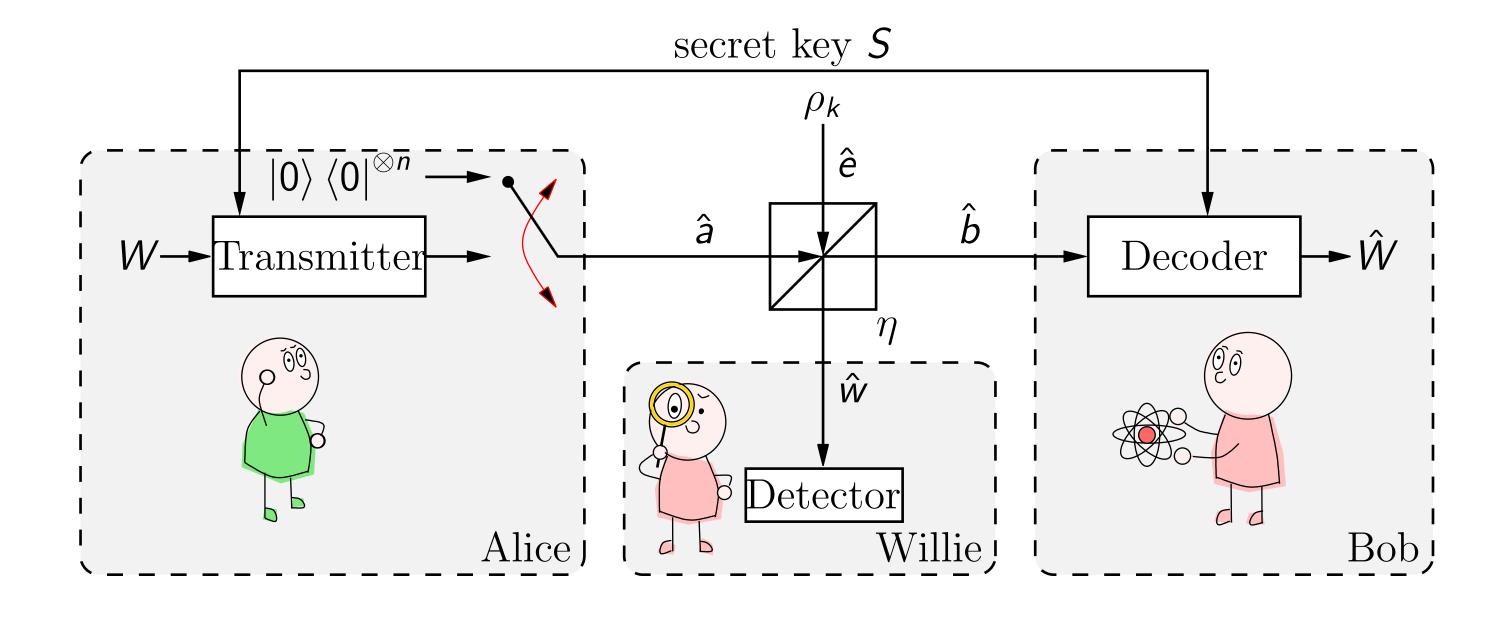
- First- and second-order asymptotics [Wang et al.'16, Bloch'16, Tahmasbi-Bloch'19]
- Multiuser channels [Arumugam-Bloch'18'19'19, Tan-Lee'19, Cho-Li'21]
- Codes for covert communications [Frèche et al.'17, Kadampot et al.'18'19, Lamarca-Matas'19, Zhang et al.'20, Wang-Bloch'21]
- AWGN channels [Wang et al.'16, Zhang et al.'19, Yan et al.'19]
- ► MIMO-AWGN channels [Abdelaziz-Koksal'17, Bendary et al.'19, Wang-Bloch'21]
- Variational distance constraint [Tahmasbi-Bloch'19, Zhang et al.'19, Wang-Bloch'21]
- Classical-Quantum channels [Bash et al.'15, Sheikholeslami et al.'16, Wang'16]
- Bosonic Channels [Bullock et al.'20, Gagatsos et al.'20]



- Switch $\phi \in \{0, 1\}$ controls transmission state at Alice
- Innocent symbol x_0 : absence of communications
 - Induces distributions $P_0 \triangleq W_{Y|X=x_0}$ and $Q_0 \triangleq W_{Z|X=x_0}$
- ightharpoonup Code C: occurrence of communications
 - induces distribution \widehat{Q}^n at Willie
- RELIABILITY: Bob reliably recovers the message and guesses when the communication occurs
- ► **DETECTION:** Willie (passive warden) distinguishes
 - Hypothesis H_0 : $Q_0^{\otimes n}$ (absence of communications)
 - ► Hypothesis H_1 : \widehat{Q}^n (occurrence of communications)
- ► GOAL: make sure Willie's test close to blind test

- Single-mode lossy thermal-noise bosonic channel $\mathcal{L}_{A \to BW}^{(\eta,k)}$
 - ightharpoonup Transmissivity η
 - Described by a beamsplitter $\hat{b}=\sqrt{\eta}\hat{a}+\sqrt{1-\eta}\hat{e}$, and $\hat{w}=\sqrt{1-\eta}\hat{a}+\sqrt{\eta}\hat{e}$
 - Environment in thermal bath with mean photon number *k*
 - Thermal state

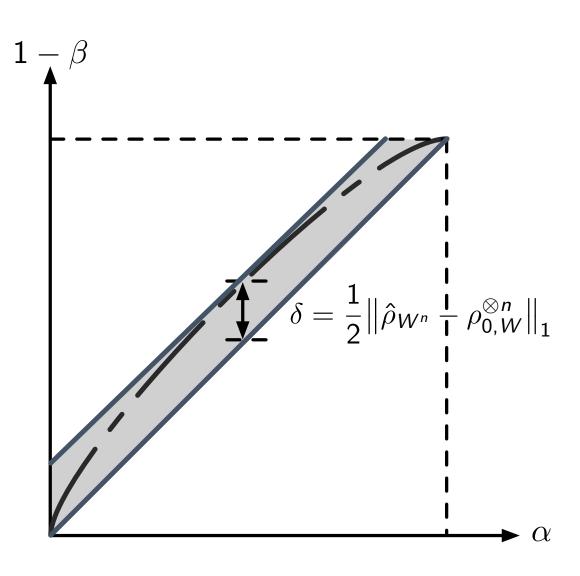
$$ho_k = rac{1}{\pi k} \int \exp\left(-rac{\left|lpha
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- ▶ **DEFINITION:** $(M, K, n, \epsilon, \delta)$ -code C
 - ▶ Uniformly distributed message $W \in [1, M]$
 - ▶ Uniformly distributed secret key $S \in [1, K]$ shared with Bob
 - ▶ Encoding channel $\mathcal{E}_{SW\to A^n}$: $|s\rangle\langle s|_S\otimes |w\rangle\langle w|_W\mapsto \rho_{A^n}(s,w)$
 - Collection of decoding POVMs $\{\{\Pi_{B^n}^{(s,w)}\}_{w \in [\![1,M]\!]}\}_{s \in [\![1,K]\!]}$
 - Induces $\hat{\rho}_{W^n} \triangleq \frac{1}{MK} \sum_{s=1}^K \sum_{w=1}^M \operatorname{tr}_{B^n} \left(\left(\mathcal{L}_{A \to BW}^{(\eta,k)} \right)^{\otimes n} \rho_{A^n}(s,w) \right)$ at Willie



Maximal average probability of error $P_e \triangleq \max_{s \in \llbracket 1, K \rrbracket} \mathbb{P} \left(\widehat{W} \neq W | S = s \right) \leqslant \epsilon$



COVERTNESS METRIC

For any test $\mathcal{T}_{W^n o \{0,1\}}$ conducted by Willie, the probabilities of false alarm α and missed detection β satisfy

$$1 \geqslant \alpha + \beta \geqslant 1 - \frac{1}{2} \|\hat{\rho}_{W^n} - \rho_{0,W}^{\otimes n}\|_1 \geqslant 1 - \delta$$

- Any test conducted by Willie is close to blind test
- Trace distance is the covertness metric that carries operational meaning

- [Tahmasbi-Bloch'19] investigates the impact and operational meaning of different covertness metrics for classical case
 - ► Relative entropy $\mathbb{D}(\widehat{Q}^n || Q_0^{\otimes n})$
 - Variational distance $\mathbb{V}(\widehat{Q}^n, Q_0^{\otimes n})$
 - Optimal probability of missed detection $\beta_{\alpha}(Q_0^{\otimes n}, \widehat{Q}^n)$

CHOICE OF COVERTNESS METRIC MATTERS

- No strong converse for covertness in covert communication
- Asymptotics of covert capacity depend on covertness metric

DRAWBACKS OF RELATIVE ENTROPY

- Loose proxy for variational distance, more stringent
- Operational meaning preserved in variational distance
- Using variational distance leads to a 25% relative increase in throughput [Wang-Bloch'21]
- TAKE AWAY: variational distance is more operationally relevant, and so is trace distance

KNOWN RELATIONS [Hayashi'17]

$$1 - \sqrt{F(\rho, \sigma)} \leqslant \frac{1}{2} \|\rho - \sigma\|_{1} \leqslant \sqrt{1 - F(\rho, \sigma)} \leqslant \sqrt{\mathbb{D}(\rho \| \sigma)}$$

- Fidelity is multiplicative for tensor-product states
- Fidelity retains properties of distance through purified distance $P(\rho, \sigma) \triangleq \sqrt{1 F(\rho, \sigma)}$
- ► [Bullock et al.'20], [Gagatsos et al.'20] use quantum relative entropy
- We use purified distance (and therefore fidelity) as intermediate metric in achievability
 - ► TECHNICAL DETAIL: require triangle inequality and purified distance for resolvability analysis

DEFINITION: ACHIEVABLE THROUGHPUT AND COVERT CAPACITY

 $(M, K, n, \epsilon, \delta)$ - code is ϵ -reliable and δ -covertness if $P_e \leqslant \epsilon$ and $\frac{1}{2} \|\hat{\rho}_{W^n} - \rho_{0,W}^{\otimes n}\|_1 \leqslant \delta$. The maximum number of messages that can be transmitted by an $(M, K, n, \epsilon, \delta)$ -code is $M^*(n, \epsilon, \delta)$.

Covert Capacity:

$$C_{\text{cov}} \triangleq \lim_{n \to \infty} \frac{\log M^*(n, \epsilon, \delta)}{\sqrt{n}}$$

For a sequence of code achieving covert capacity, secret key throughput is $\lim_{n\to\infty} \frac{\log K}{\sqrt{n}}$.

THEOREM: COVERT CAPACITY OF LOSSY THERMAL-NOISE BOSONIC CHANNEL

Covert capacity satisfies

$$\frac{2\sqrt{\eta k(\eta k+1)}}{1-\eta}\eta\log\left(1+\frac{1}{(1-\eta)k}\right)Q^{-1}\left(\frac{1-\delta}{2}\right)\geqslant C_{\text{cov}}\geqslant\frac{2\sqrt{\eta k(\eta k+1)}}{1-\eta}\eta\log\left(1+\frac{1}{(1-\eta)k}\right)\delta$$

Lower bound is achievable with key throughput satisfying

$$2\sqrt{\eta k(\eta k+1)} \left(\log \left(1 + \frac{1}{\eta k} \right) - \frac{\eta}{1-\eta} \log \left(1 + \frac{1}{(1-\eta)k} \right) \right)^{+} \delta$$

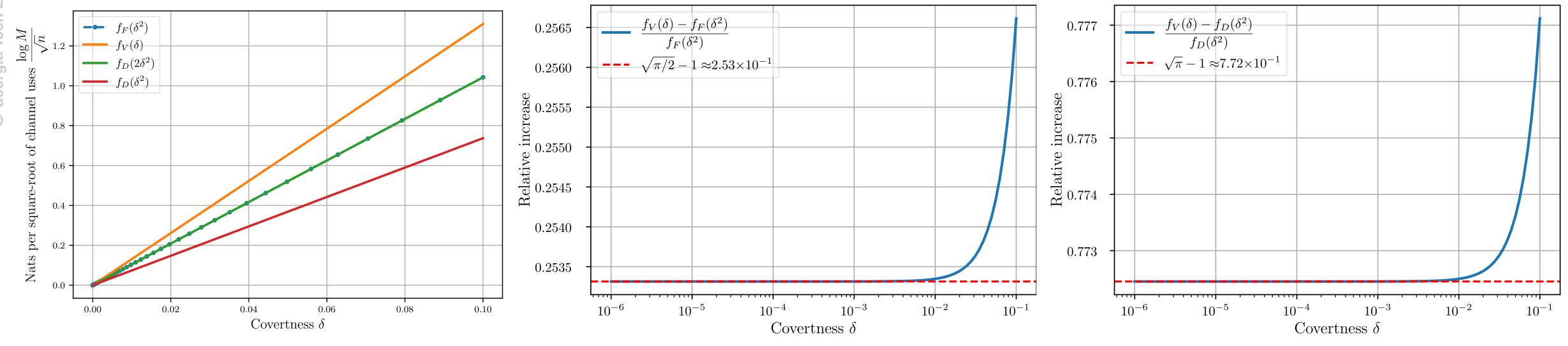
where $(x)^+ \triangleq \max(x, 0)$.

TAKE AWAY: covert capacity is a function of covertness metric.

- ► **COMPARISON:** given $\frac{1}{2} \|\hat{\rho}_{W^n} \rho_{0,W}^{\otimes n}\|_1 \leqslant \sqrt{1 F(\hat{\rho}_{W^n}, \rho_{0,W}^{\otimes n})} \leqslant \sqrt{\mathbb{D}(\hat{\rho}_{W^n} \| \rho_{0,W}^{\otimes n})} \leqslant \delta$
 - Upper-bound is derived with trace distance metric, denoted by $f_V(\delta)$
 - Lower-bound is derived with fidelity metric, denoted by $f_F(\delta^2)$
 - We also compare previous result by [Gagatsos et al.'20] with quantum relative entropy metric, denoted by $f_D(\delta^2)$
- ► However, [Bullock et al.'20] and [Gagatsos et al.'20] consider quantum Pinsker's inequality
 - We denote this result by $f_D(2\delta^2)$

$$\frac{1}{2} \|\hat{\rho}_{W^n} - \rho_{0,W}^{\otimes n}\|_1 \leqslant \sqrt{\frac{\mathbb{D}(\hat{\rho}_{W^n} \| \rho_{0,W}^{\otimes n})}{2}} \leqslant \delta$$

From above inequalities, $f_V(\delta) \geqslant f_F(\delta^2) = f_D(2\delta^2) \geqslant f_D(\delta^2)$



- ▶ **GOAL:** use **fidelity** to ensure $\frac{1}{2} \|\hat{\rho}_{W^n} \rho_{0,W}^{\otimes n}\|_1 \leqslant \sqrt{1 F(\hat{\rho}_{W^n}, \rho_{0,W}^{\otimes n})} \leqslant \delta$
- Coherent-state QPSK signaling with mean photon number decreasing in $\mathcal{O}\left(\frac{1}{\sqrt{n}}\right)$
 - Ensure signal to be close to innocent states
 - Precise control of covertness
- One-shot channel reliability and resolvability
 - Existence of codebook that is reliable and does not change covertness much
 - Second-order asymptotic

- Coherent-state QPSK with **decreasing** mean photon number Generated state at Alice's system $\rho_{n,XA} \triangleq \sum_{x \in \llbracket 0,3 \rrbracket} \frac{1}{4} |x\rangle \langle x|_X \otimes \left| u_n e^{j\pi x/2} \right\rangle \langle u_n e^{j\pi x/2} |_A$ Induced state at Willie's system $\rho_{n,XW} \triangleq \sum_{x \in \llbracket 0,3 \rrbracket} \frac{1}{4} |x\rangle \langle x|_X \otimes \rho_{\eta k,W} (\sqrt{1-\eta} u_n e^{j\pi x/2})$

 - **INTUITION:** perturbation in displacement around thermal state
- Perturbation theory for quantum fidelity [Grace-Guha'21]

EMMA: FIDELITY ANALYSIS OF COHERENT-STATE QPSK By setting
$$u_n \triangleq \frac{\left(4\eta k(\eta k+1)(\delta-2\kappa)^2\right)^{\frac{1}{4}}}{\sqrt{1-\eta}n^{\frac{1}{4}}}$$
, we have
$$F\left(\rho_{n,W}^{\otimes n},\rho_{0,W}^{\otimes n}\right) \geqslant 1-(\delta-2\kappa)^2+\frac{(\delta-2\kappa)^4}{2}.$$

Make sure existence of a codebook that is both reliable and covert

INGREDIENT

- Position-based coding and sequential decoding [Wilde'17], [Oskouei et al.'18]
- Convex-split lemma [Anshu et al.'19], [Khatri et al.'19]

LEMMA: ONE-SHOT CHANNEL RELIABILITY AND RESOLVABILITY

By choosing fix ϵ , κ , γ_1 , γ_2 , $\gamma_3 > 0$, then for a bi-partite state ρ_{XA} and a channel

 $\mathcal{G}: \rho_{XA} \mapsto \rho_{XBW}$, there exists a coding scheme such that

$$\log M \geqslant \mathbb{D}_{H}^{\epsilon^{2}/10-\gamma_{1}}(\rho_{XB} \| \rho_{X} \otimes \rho_{B}) - \log\left(\frac{4\epsilon^{2}}{10\gamma_{1}^{2}}\right),$$

$$\log MK \leqslant \mathbb{D}_{\max}^{\kappa/2-\gamma_{2}-\gamma_{3}}(\rho_{XW} \| \rho_{X} \otimes \rho_{W}) + 2\log\left(\frac{1}{\gamma_{2}}\right) + \log\left(\frac{8}{\gamma_{3}^{2}}\right),$$

$$\mathbb{E}_{\mathcal{C},S}\Big\{\mathbb{P}\Big(\widehat{W} \neq W|S\Big)\Big\} \leqslant \frac{\epsilon^{2}}{10},$$

$$\mathbb{E}_{\mathcal{C}}\Big\{\frac{1}{2}\|\widehat{\rho}_{W} - \rho_{W}\|_{1}\Big\} \leqslant \kappa - \gamma_{2},$$

where \mathcal{C} is the codebook.

- Analyze trace distance directly with techniques developed in [Tahmasbi-Bloch'19, Zhang et al.'19, Wang-Bloch'21]
- Constrained to coherent-state codebook
 - A weaker converse
 - No entanglement within modes
 - Possible to lift this constraint
- Establish lower bound on trace distance related to minimum received photon number of codewords
 - Require us to find a simple to analyze yet powerful test for Willie
 - Photon counter
- Show existence of good sub-code resulting in low covertness metric
- Obtain upper bound on covert message size of good sub-code

Suboptimal photon counting test POVM $\{T, I - T\}$

$$T \triangleq \sum_{m_1+m_2+\cdots+m_n \geqslant \tau} |m_1\rangle \langle m_2| \otimes |m_2\rangle \langle m_2| \otimes \cdots \otimes |m_n\rangle \langle m_n|$$

- Since modes are in tensor-product states, it reduces to a classical test $T(m^n) = \mathbb{I}\left\{\sum_{i=1}^n m_i \geqslant \tau\right\}$
- Carefully choose threshold and use photon number statistics of displaced thermal states with Berry-Esseen Theorem

LEMMA: LOWER BOUND ON TRACE DISTANCE

By setting
$$au = \frac{(1-\eta)N_*}{2} + n\eta k$$
, we have
$$\frac{1}{2} \|\hat{\rho}_{W^n} - \rho_{0,W}^{\otimes n}\|_1 \geqslant 1 - \alpha - \beta$$

$$\geqslant 1 - 2Q \left(\frac{(1-\eta)N_*}{2\sqrt{n\eta k(\eta k+1)}} \right) - \frac{(1-\eta)^2(1+2\eta k)N_*^2}{4\sqrt{2\pi}n^{3/2}[\eta k(\eta k+1)]^{3/2}} - \frac{B_0 + B_1}{\sqrt{n}},$$

where $N_* \triangleq \min_{m \in \mathcal{M}} \sum_{i=1}^{n} \left| \alpha_i^{(m)} \right|^2$ is minimum photon number of codewords, α is false-alarmed probability, β is missed-detection probability, and B_0 , B_1 are some constants.

- ► INTUITION: covertness metric is higher in a set of high-photon-number codewords
- Partition code $\mathcal C$ into high-photon-number (bad) sub-code $\mathcal C^{(h)}$ and low-photon-number (good) sub-code $\mathcal C^{(\ell)}$

LEMMA: EXISTENCE OF GOOD SUB-CODE

For any covert codebook \mathcal{C} , let $\lim_{n\to\infty}\gamma_n=0$. Then there exists a sub-code $\mathcal{C}^{(\ell)}$ such that $\left|\mathcal{C}^{(\ell)}\right|\geqslant\gamma_n\left|\mathcal{C}\right|$ and photon-number of codewords $N^{(c)}\leqslant A\sqrt{n}$,

$$A \triangleq \frac{2\sqrt{\eta k(\eta k+1)}}{1-\eta}Q^{-1}\left(\frac{1-\delta}{2} - \frac{\nu^2(1-\eta)^2(1+2\eta k)}{4\sqrt{2\pi n}[\eta k(\eta k+1)]^{3/2}} - \gamma_n\right),$$

where ν depends on channel.

Characterize upper bound on photon-number of codewords in good sub-code

- ▶ Code C consists of K sub-codes C_s indexed by key $s \in [1, K]$
 - Size of each sub-code is M
 - $\mathcal{C}_s^{(\ell)} \triangleq \mathcal{C}_s \cap \mathcal{C}^{(\ell)}$
 - By pigeonhole principle, $\left|\mathcal{C}_{s}^{(\ell)}\right|\geqslant\gamma_{n}M$
- We find an upper bound on $\log \left| \mathcal{C}_s^{(\ell)} \right|$ and therefore apply it to M
 - For bosonic channel, what we need is the bound on photon number we derived earlier
 - By Fano's inequality, Holevo bound, and capacity of bosonic channel

$$\log \left| \mathcal{C}_s^{(\ell)} \right| (1 - \frac{\epsilon_n}{\gamma_n}) - 1 \leqslant nC(\frac{A}{\sqrt{n}}, \eta, k) \leqslant \eta A \sqrt{n} \log \left(1 + \frac{1}{(1 - \eta)k} \right)$$

$$\liminf_{n \to \infty} \frac{\log M}{\sqrt{n}} \leqslant \liminf_{n \to \infty} \frac{\log \left| \mathcal{C}_s^{(\ell)} \right| - \log \gamma_n}{\sqrt{n}}$$

$$= \frac{2\sqrt{\eta k(\eta k + 1)}}{1 - \eta} \eta \log \left(1 + \frac{1}{(1 - \eta)k} \right) Q^{-1} \left(\frac{1 - \delta}{2} \right)$$

 $ightharpoonup C(N, \bar{N}, \eta)$ is channel capacity of bosonic channel with photon numbers of two ports N, \bar{N} and transmissivity η

- Solution Gaussian ensemble $\left\{ \frac{1}{\pi s_n} \exp\left(-\frac{|lpha|^2}{s_n}\right)$, $|lpha\rangle\langlelpha| \right\}$ with $s_n = \frac{2\sqrt{\eta k(\eta k+1)}}{1-\eta}Q^{-1}\left(\frac{1-\delta}{2}\right)$ would achieve upper-bound solution.
 - ► TECHNICAL DIFFICULTY: our coding theorem does not support uncountable alphabet
- Direct analysis of QPSK signaling with trace distance is challenging
- To extend converse to general *n*-mode states, consider applying an entanglement-breaking processing channel

$$\rho_W \mapsto \sum_{m} \operatorname{tr}(|m\rangle\langle m|\,\rho_W) \otimes |m\rangle\langle m|_M$$

- Destroy entanglement between modes
- Any single-mode Gaussian state is thermal state with displacement and symplectic transform

CONCLUSION

- Covert capacity of bosonic channel depends on choice of covertness metric
- Trace distance is ultimate covertness metric carrying operational meaning
- Characterize secret-key throughput of bosonic covert communications



Fundamental Limits of Quantum-Secure Covert Communication Over Bosonic Channels

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Covert Capacity of Bosonic Channels

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Covert MIMO Communications under Variational Distance Constraint

S.-Y. Wang and M. R. Bloch

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Position-based Coding and Convex Splitting for Private Communication over Quantum Channels

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Quantum Information Processing, vol. 16, no. 10, p. 264, Oct. 2017



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